Algorithms for solving rational quadratic forms *

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Abstract

We present algorithms for solving quadratic forms over rational numbers based on constructive proofs of the Hasse principle for these forms.

1 Introduction

In this document we describe algorithms for finding a nontrivial rational solution of a given quadratic form with any number of variables.

Quadratic forms (i.e. homogenous quadratic polynomials) over rationals belong to the class of polynomials where the so-called Hasse principle holds. Namely, there exists a rational solution to the form if and only if there exists a solution in every local field containing \mathbb{Q} , it means any field complete with respect to some valuation defined on \mathbb{Q} .

In the case of quadratic forms there is even a stronger result: we can construct a global (means rational) solution of a given quadratic form by combining local solutions (means solutions in local fields).

For developing the algorithms we mainly follow [2]. After implementing the algorithm for solving ternary quadratic forms we learned that meanwhile there was another algorithm developed (cf. [3]), which avoids the expensive LLL reduction a is therefore faster. So in the last version of the package we replaced our version by that one. However, algorithms for solving forms with more variables are new.

It is worth to mention that the most interesting cases for solving quadratic forms are those with three and four variables. The reason is, that a form of 5 and more variables has a solution in every local field except possibly the field of real numbers (see Lemma 16). So, by the Hasse principle, if we want to test whether there is a rational solution, it is enough to test the solvability in \mathbb{R} , which is very easy.

We implemented all algorithms in Magma ([1]). They should appear in one of next versions of the system.

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2 The field of *p*-adic numbers

Definition 1. Let k be any field. The function $|.|: k \to \mathbb{R}$ is called a valuation if it satisfies the following conditions:

- (i) $|a| \ge 0 \ \forall a \in k$, and |a| = 0 iff a = 0.
- (ii) $|ab| = |a||b| \quad \forall a, b \in k.$
- (iii) $|a+b| \le |a| + |b| \ \forall a, b \in k$ (triangle inequality).

Here we consider only valuations defined for the field of rational numbers. We have the following possibilities how to define a non-trivial valuation:

- absolute value,
- *p*-adic valuation: For a fixed prime *p* let

$$r = p^{\rho} u/v; \quad u, v \in \mathbb{Z}, \ p \nmid uv.$$

Then we define a non-archimedean valuation

$$|r|_p = p^{-\rho}$$

Usually the standard absolute value is called *valuation at infinity* and denoted by $|.|_{\infty}$.

Theorem 2. (Ostrowski) Every non-trivial valuation on \mathbb{Q} is equivalent to one of the valuations $|.|_p$, where either p is a prime number or $p = \infty$.

Every valuation gives rise to a topology on \mathbb{Q} . And like the field \mathbb{R} of real numbers is constructed as the completion of the field \mathbb{Q} w.r.t. the absolute value, we have also the following

Definition 3. Let $|.|_p$ be a non-archimedean valuation on \mathbb{Q} . The field of p-adic numbers \mathbb{Q}_p is the completion of \mathbb{Q} with respect to the valuation $|.|_p$.

Again, as in the case of absolute value, one often denotes \mathbb{R} by \mathbb{Q}_{∞} , since it is the completion of \mathbb{Q} w.r.t. $|.|_{\infty}$. Traditionally, the field of rational numbers is referred to as the *global field*, and all its completions (i.e. all *p*-adic fields together with $\mathbb{R} = \mathbb{Q}_{\infty}$) are called *local fields*.

Since every local field \mathbb{Q}_p (*p* being prime or ∞) contains \mathbb{Q} as its subset, it is clear, that if a polynomial equation with rational coefficients has a solution over \mathbb{Q} , then this is also the solution over any local field. So we can decide about unsolvability of the equation over \mathbb{Q} by finding one local field, where it is not solvable.

It would be very nice, if also the converse holds, namely:

Local-Global Principle. The existence or non-existence of solutions in \mathbb{Q} (global solution) of a diophantine equation can be detected by studying the solutions of the equation in \mathbb{Q}_p , p - prime or ∞ (local solutions).

Example. Sometimes the Hasse Principle does not work:

$$(x^2 - 2)(x^2 - 17)(x^2 - 34) = 0$$

is solvable locally everywhere, but is not solvable globally, i.e. it has no rational solution.

Although the Hasse principle is not valid for all polynomials, still there are important classes of equations, for which it holds. One of such classes are quadratic forms over \mathbb{Q} . For these polynomials we have even more: we are able not only to decide the existence of a rational solution, but we can also construct the rational solution from the local ones.

3 Rational quadratic forms

3.1 Basic notions

Definition 4. Let k be a field of characteristic different from 2. By n-ary quadratic form over k we understand a homogeneous quadratic polynomial over k with n variables.

Using a linear transformation of variables any quadratic form can be brought into a diagonal form:

$$f(\mathbf{x}) = a_1 x_1^2 + a_2 x_2^2 + \dots + a_n x_n^2; \quad a_i \in k.$$
(1)

So from now on we suppose that the quadratic form is given as in (1).

Definition 5. The form (1) is regular, if $a_i \neq 0 \forall i$.

A regular quadratic form is said to be isotropic, if there are $b_1, \ldots, b_n \in k$ not all 0, such that $a_1b_1^2 + \cdots + a_nb_n^2 = 0$. The vector (b_1, \ldots, b_n) is then called an isotropic vector.

The form (1) represents $c \in k$, if there are $c_1, \ldots, c_n \in k$ such that $a_1c_1^2 + \cdots + a_nc_n^2 = c$. The form over k is universal if it represents every non-zero element of k.

Lemma 6. Let $f(\mathbf{x}) = a_1 x_1^2 + a_2 x_2^2$ be isotropic. Then it is universal.

3.2 Hasse principle for ternary quadratic forms

Using a linear transformation of variables we can any ternary quadratic form bring into the shape

 $f(\mathbf{x}) = a_1 x_1^2 + a_2 x_2^2 + a_3 x_3^2$ with $a_i \in \mathbb{Z}$ and $a_1 a_2 a_3$ square-free.

To use the Hasse principle for deciding solvability and solving the form over rational numbers, we first have to solve the problem for infinitely many local fields. But the Lemma 7 allows us to restrict the decision process to finitely many cases. Afterwards the Lemma 8 says us that solving the quadratic form in a local field is equivalent to solving the form modulo prime, resp. power of prime.

Lemma 7. Let $p \neq 2, \infty$ and let

$$f(\mathbf{x}) = a_1 x_1^2 + a_2 x_2^2 + a_3 x_3^2$$
 with $|a_1|_p = |a_2|_p = |a_3|_p$.

Then $f(\mathbf{x})$ is isotropic over \mathbb{Q}_p .

So we have to test the solvability only in \mathbb{R} , \mathbb{Q}_2 and those \mathbb{Q}_p , where p divides one of coefficients a_1, a_2, a_3 .

Lemma 8. Let

$$f(\mathbf{x}) = a_1 x_1^2 + a_2 x_2^2 + a_3 x_3^2$$

with $a_1, a_2, a_3 \in \mathbb{Z}$ and $a_1a_2a_3$ square-free.

- (i) Let p be an odd prime, $p \mid a_3$. The form $f(\mathbf{x})$ is isotropic in \mathbb{Q}_p if and only if there is an $r \in \mathbb{Z}$ such that $a_1r^2 + a_2 \equiv 0 \pmod{p}$.
- (ii) Let $2 \mid a_3$. The form $f(\mathbf{x})$ is isotropic in \mathbb{Q}_2 if and only if there is an s = 0 or 1 such that

$$a_1 + a_2 + a_3 s^2 \equiv 0 \pmod{8}.$$

(iii) Let $2 \nmid a_1 a_2 a_3$. The form $f(\mathbf{x})$ is isotropic in \mathbb{Q}_2 if and only if after suitable permuting the suffices 1, 2, 3 we have

$$a_1 + a_2 \equiv 0 \pmod{4}.$$

Theorem 9. (Legendre) Let

$$f(\mathbf{x}) = a_1 x_1^2 + a_2 x_2^2 + a_3 x_3^2$$

with $a_1, a_2, a_3 \in \mathbb{Z}$ and $a_1a_2a_3$ square-free. Suppose that the following conditions are satisfied

(i) if p is an odd prime dividing $a_1a_2a_3$, say $p \mid a_3$, then there is an integer r_p such that

$$a_1 r_p^2 + a_2 \equiv 0 \pmod{p}.$$

(*ii*) if $2 | a_1 a_2 a_3$, say $2 | a_3$ then

$$a_1 + a_2 + a_3 s^2 \equiv 0 \pmod{8}$$

where s = 0 or 1.

(iii) if $2 \nmid a_1 a_2 a_3$ then, on permuting a_1, a_2, a_3 , if need be,

$$a_1 + a_2 \equiv 0 \pmod{4}.$$

Then there are $b_1, b_2, b_3 \in \mathbb{Z}$, not all 0, such that

$$a_1b_1^2 + a_2b_2^2 + a_3b_3^2 = 0.$$

A sketch of proof: The idea of the proof is, that the modular quadratic equation can be reformulated into a linear condition for the isotropic vector.

(i) Let p be an odd prime, $p \mid a_3$. The condition from the Theorem can be written as the linear condition for the solution, namely

$$z_1 \equiv r_p z_2 \pmod{p}$$
.

Then $f(\mathbf{z}) \equiv 0 \pmod{p}$.

(ii) Let $2 \mid a_3$. Then the linear conditions

 $z_1 \equiv z_2 \pmod{4}$ $z_3 \equiv sz_2 \pmod{2}$

imply that $f(\mathbf{z}) \equiv 0 \pmod{8}$.

(iii) Let $2 \nmid a_1 a_2 a_3$. Then the linear conditions

 $z_1 \equiv z_2 \pmod{2}$ $z_3 \equiv 0 \pmod{2}$

imply that $f(\mathbf{z}) \equiv 0 \pmod{4}$.

So we have a collection of linear modular conditions for the isotropic vector and using Chinese remainder theorem we combine them into one linear modular equation, such that every its solution \mathbf{z} is also a solution to the modular quadratic equation

$$f(\mathbf{z}) \equiv 0 \pmod{4a_1a_2a_3}.$$

The solutions to the linear equation form a lattice in \mathbb{Z}^3 . By the results from geometry of numbers we know, that among the points in the lattice there exists also a point, which is a solution to the quadratic form over rational numbers, for example the shortest vector in the lattice. Therefore we finally use LLL-reduction to find the solution. \Box

Remark. There are two expensive actions in the algorithm: first one is the factorization of the coefficients, and the second one is the final LLL-reduction. After implementing the algorithm we learned, that there is already an algorithm avoiding the LLL-reduction implemented in Magma, see [3]. Therefore at the end we replaced our LLL-based algorithm by the more efficient one.

3.3 Hasse principle for quaternary quadratic forms

While solving ternary quadratic forms we could avoid the work with p-adic numbers. First, finding a p-adic solution was equivalent to finding a solution modulo p for odd prime p, or modulo 4 or 8 for p = 2. And afterwards instead of combining particular p-adic solutions into one rational solution of the quadratic form, it was enough to combine modular solutions of linear equations via Chinese remainder theorem.

However, for solving quaternary forms, it is not possible anymore to escape from the local fields. For this, let's first have a closer look at the solvability of ternary quadratic forms over local fields.

For the multiplicative group \mathbb{Q}_p^* of every local field, the subset $(\mathbb{Q}_p^*)^2$ is its subgroup. Therefore we can form a factor group $\mathbb{Q}_p^*/(\mathbb{Q}_p^*)^2$. One can prove that $\mathbb{Q}_p^*/(\mathbb{Q}_p^*)^2$ is always of finite order. For example, for $\mathbb{Q}_{\infty}^* = \mathbb{R}^*$ we have $(\mathbb{R}^*)^2 = \mathbb{R}^+$, the group of all positive real numbers. So $|\mathbb{R}^*/(\mathbb{R}^*)^2| = 2$. If p is an odd prime, then $|\mathbb{Q}_p^*/(\mathbb{Q}_p^*)^2| = 4$ and for p = 2, we have $|\mathbb{Q}_2^*/(\mathbb{Q}_2^*)^2| = 8$.

Definition 10. Hilbert norm residue symbol *is defined for* $a, b \in \mathbb{Q}_p^*$ *as follows:* $(.,.): \mathbb{Q}_p^* \times \mathbb{Q}_p^* \to \{1,-1\},$

$$(a,b) = \begin{pmatrix} a,b\\ p \end{pmatrix} = \begin{cases} 1 & \text{if } ax^2 + by^2 - z^2 \text{ is isotropic} \\ -1 & \text{otherwise.} \end{cases}$$

The Hilbert norm residue symbol (a, b) depends only on a, b modulo squares: $(a, b) = (ac^2, bd^2)$ for any $c, d \in \mathbb{Q}_p^*$, since any (x_0, y_0, z_0) is a solution to $ax^2 + by^2 - z^2$ if and only if $(x_0/c, y_0/d, z_0)$ is a solution to $(ac^2)x^2 + (bd^2)y^2 - z^2$. Therefore (., .) is well defined on quadratic residuum classes. It follows that when deciding solvability of a ternary quadratic form over a local field, we can reduce to finitely many cases. We just have to find the residuum classes of the coefficients of the quadratic form and then look up the value of the Norm residue symbol in the table.

Theorem 11. (Product formula) Let $a, b, \in \mathbb{Q}^*$. Then $\binom{a,b}{p} = 1$ for almost all p and

$$\prod_{p \le \infty} \binom{a, b}{p} = 1$$

Remark. The form $f(\mathbf{x}) = a_1 x_1^2 + a_2 x_2^2 + a_3 x_3^2$ is isotropic over \mathbb{Q}_p if

$$\binom{-a_1a_3, -a_2a_3}{p} = 1.$$

By the Product formula

$$\prod_{p \le \infty} \binom{-a_1 a_3, -a_2 a_3}{p} = 1.$$

So while testing solvability over local fields, we can omit one of them. This also explains, why in the Legendre Theorem the solvability of the form over \mathbb{R} need not to be tested.

For solving a quaternary quadratic form

$$f(\mathbf{x}) = a_1 x_1^2 + a_2 x_2^2 + a_3 x_3^2 + a_4 x_4^2$$
 with $a_j \in \mathbb{Z}, \ j = 1, \dots, 4,$

we will split it into two ternary forms as follows:

$$a_1 x_1^2 + a_2 x_2^2 - t y_1^2 a_3 x_3^2 + a_4 x_4^2 + t y_2^2$$

for suitable $t \in \mathbb{Z}^*$. If the two ternary forms are isotropic, then apparently also the quaternary form is isotropic and we can construct its isotropic vector from the ones of two ternary forms and vice versa: isotropy of the quaternary forms enables us to construct the two ternary isotropic forms. By Lemma 6 we can in this case always construct regular ternary forms, i.e. with $t \neq 0$. So in our algorithm we will reduce the problem of solving a quaternary form to the problem of constructing the $t \in \mathbb{Z}$. We will do it in two steps: first we construct such t_p in every relevant local field \mathbb{Q}_p and afterwards we combine the "local" t_p 's into a "global" $t \in \mathbb{Q}$. The following Lemmas explain how to find a t_p in a local field.

Lemma 12. Let $g(\mathbf{x}) = a_1 x_1^2 + a_2 x_2^2$ be a regular form over Q_p . Then $b \neq 0$ is represented by g if and only if

$$(b, -a_1a_2) = (a_1, a_2).$$

Lemma 13. The quaternary form $f(\mathbf{x}) = a_1 x_1^2 + a_2 x_2^2 + a_3 x_4^2 + a_4 x_4^2$ is anisotropic if and only if

$$\begin{aligned} -a_1 a_2(\mathbb{Q}_p^*)^2 &= -a_3 a_4(\mathbb{Q}_p^*)^2 \quad and \\ (a_1, a_2) &= -(-a_3, -a_4). \end{aligned}$$

The Corollary of the following Theorem plays the same role in solving quaternary forms, as the Chinese remainder theorem in solving ternary forms: It combines local "solutions" $(t_p \in \mathbb{Q}_p)$ into a global "solution" $(t \in \mathbb{Q})$.

Theorem 14. (Dirichlet) Let l > 0 and m be coprime pair of integers. There are infinitely many primes $p \equiv m \pmod{l}$.

Corollary 15. Let P be finite set of primes p (possibly also ∞). For each $p \in P$ let $t_p \in \mathbb{Q}_p^*$ be given. Then there is a $t \in \mathbb{Q}^*$ such that

- (i) $t \in t_p(\mathbb{Q}_p^*)^2$ for all $p \in P$,
- (ii) $|t|_p = 1$ for all $p \notin P, p \neq \infty$ except possibly for one $p = p_0$.

Proof: $p \in P, p \neq \infty$: $t_p = p^{\alpha(p)}s_p$, $|s_p|_p = 1$. We look for t in the form:

$$t = \pm p_0 \prod_{p \in P \setminus \{\infty\}} p^{\alpha(p)}, \quad p_0 \notin P \text{ prime.}$$

The sign is chosen according to p_{∞} (if $\infty \in P$). According to the Dirichlet theorem we can choose prime $p_0 > 0$ such that $\forall p \in P, p \neq \infty$

$$p^{-\alpha(p)}t \equiv s_p \begin{cases} \mod p & p \neq 2\\ \mod 8 & p = 2 \end{cases}$$

Then the t satisfies the desired properties. \Box

Now we have constructed $t \in \mathbb{Q}$, such that the forms

$$a_1x_1^2 + a_2x_2^2 - ty_1^2, \qquad a_3x_3^2 + a_4x_4^2 + ty_2^2 \tag{2}$$

are isotropic over \mathbb{Q}_p , if the forms

$$a_1x_1^2 + a_2x_2^2 - t_py_1^2$$
, $a_3x_3^2 + a_4x_4^2 + t_py_2^2$

were isotropic over \mathbb{Q}_p . This holds for $p = 2, \infty$ and every prime dividing one of coefficients of the quaternary form (follows from the part (i) of the Corollary) as well as any prime not involved in the coefficients (follows from the part (ii)), up to possibly one prime p_0 . Therefore the two ternary forms (2) are isotropic over every $\mathbb{Q}_p, p \neq p_0$. From the product formula (Theorem 11) we have, that they are isotropic over \mathbb{Q}_{p_0} too, and hence by the Hasse principle isotropic over \mathbb{Q} .

3.4 Quadratic forms with 5 and more variables

Lemma 16. Let $n \ge 5$, $p < \infty$. A form in n variables over \mathbb{Q}_p is isotropic.

Since the Hasse principle holds for all quadratic forms (see [2]), for forms with at least 5 variables we have a very easy criterion for solvability: A form is isotropic over \mathbb{Q} if and only if it is isotropic over \mathbb{R} .

4 The Algorithms

4.1 Isotropic vector of ternary quadratic form ¹

FUNCTION SolubilityCertificate

INPUT: Integers b_1, b_2, b_3 , where all b's are square-free and pairwise coprime, and b_3 is odd.

OUTPUT: Integer r such that $r^2 = -b_1/b_2 \pmod{b_3}$ if exists.

 $^{^1\}mathrm{The}$ algorithm for ternary forms is obsolete. Instead, the algorithm described in [3] is used.

1. for each prime p dividing b_3

if $-b_1/b_2 \pmod{p}$ is a perfect square over Z_p then $r_p := \sqrt{-b_1/b_2}$ (in Z_p) else print "not solvable"; exit end if end for;

- 2. By Chinese reminder theorem combine r_p 's (i.e. the solutions of the system $r_p = \sqrt{-b_1/b_2}$ (in Z_p)) to r (i.e. the solution of $r = \sqrt{-b_1/b_2}$ (in Z_{b_3})).
- 3. return r.

FUNCTION IsIsotropic3

INPUT: Coefficients c_1, c_2, c_3 of a diagonal form $c_1x_1^2 + c_2x_2^2 + c_3x_3^2$. OUTPUT: **true** if and only if the form is isotropic, and if so, returns an isotropic vector.

- 1. Solve trivial case, when one of the coefficients is zero.
- 2. Simple test of solvability: coefficients should be not all of the same sign (\sim solvability over R).
- 3. Transform the form $c_1x_1^2 + c_2x_2^2 + c_3x_3^2$ to the form $a_1x_1^2 + a_2x_2^2 + a_3x_3^2$ so that $a_1a_2a_3$ is square-free and create the matrix T transforming solutions of $a_1x_1^2 + a_2x_2^2 + a_3x_3^2$ to the solutions of $c_1x_1^2 + c_2x_2^2 + c_3x_3^2$.
- 4. Compute the solubility certificate (exclude prime 2):

 $r_3 :=$ SolubilityCerificate $(a_1, a_2, \tilde{a}_3);$ $r_1 :=$ SolubilityCerificate $(a_2, a_3, \tilde{a}_1);$

 $r_2 :=$ SolubilityCerificate $(a_3, a_1, \tilde{a}_2);$

where \tilde{a}_i equals a_i if a_i is odd, and $a_i/2$ otherwise.

5. By Chinese reminder theorem combine the three modular equations

$$\begin{array}{rrrr} r_3 x_1 - & x_2 &= 0 \pmod{\tilde{a}_3} \\ r_1 x_2 - & x_3 = 0 \pmod{\tilde{a}_1} \\ -x_1 & +r_2 x_3 = 0 \pmod{\tilde{a}_3} \end{array}$$

into one equation

$$k_1x_1 + k_2x_2 + k_3x_3 = 0 \pmod{\tilde{a}_1\tilde{a}_2\tilde{a}_3}.$$

- 6. Add two extra modular equations for prime 2 (mod 2,4).
- 7. Compute the solution lattice of the three equations (modulo $4a_1a_2a_3$).
- 8. Find the shortest vector in the lattice.
- 9. Using the matrix T transform the solution of $a_1x_1^2 + a_2x_2^2 + a_3x_3^2$ to the solution of $c_1x_1^2 + c_2x_2^2 + c_3x_3^2$.

4.2 Isotropic vector of quaternary quadratic form

FUNCTION: IsIsotropic4

INPUT: Coefficients a_1, a_2, a_3, a_4 of a diagonal form $a_1x_1^2 + a_2x_2^2 + a_3x_3^2 + a_4x_4^2$. OUTPUT: **true** if and only if the form is isotropic, and if so, returns an isotropic vector.

The algorithm constructs such t that there are b_1, b_2, b_3, b_4 :

$$t = a_1 b_1^2 + a_2 b_2^2,$$

$$t = -a_3 b_3^2 - a_4 b_4^2$$

Afterwards, two ternary forms

$$a_1x_1^2 + a_2x_2^2 - tu^2 = 0$$

$$a_3x_3^2 + a_4x_4^2 + tv^2 = 0$$

are solved.

- 1. Solve trivial case, when one of the coefficients is zero.
- 2. Create the set P of all relevant primes; $P = \{p; p | 2a_1a_2a_3a_4\} \cup \{\infty\}$.
- 3. For each prime p in the set P find t_p (resp. its residue class in $Q_p^\ast|(Q_p^\ast)^2)$ such that

$$t_p = a_1 b_{1p}^2 + a_2 b_{2p}^2 \qquad \text{for some } b_{1p}, b_{2p} \in Q_p \\ t_p = -a_3 b_{3p}^2 - a_4 b_{4p}^2 \qquad \text{for some } b_{3p}, b_{4p} \in Q_p.$$

Such t_p exists iff (Lemma 13)

$$(a_1, a_2) = (-a_3, -a_4)$$
 or $-a_1 a_2 (Q_p^*)^2 \neq -a_3 a_4 (Q_p^*)^2$.

Otherwise print "No solution" and stop. (Both, decision of existence and finding the residue class for t_p is made by looking up in the table.)

4. Construct

$$q = \pm \prod_{p \in P \setminus \{\infty\}} p^{\beta(p)} \text{ with } \beta(p) = \begin{cases} 1 & \text{if in } |t_p|_p = p^{-\alpha(p)} \text{ is } \alpha(p) \text{ odd,} \\ 0 & \text{otherwise.} \end{cases}$$

The sign is chosen according to t_{∞} .

(Note: This and the previous step are done simultaneously.)

5. Generate all primes p_0 starting from 2. For each p_0 test if for $t = p_0 q$ are forms

$$a_1x_1^2 + a_2x_2^2 - tu^2 = 0$$

$$a_3x_3^2 + a_4x_4^2 + tv^2 = 0$$

isotropic, i.e. if for each prime $p \in P - \{\infty\}$

$$(t, -a_1a_2) = (a_1, a_2)$$

 $(t, -a_3a_4) = (-a_3, -a_4).$

6. Solve ternary forms

$$a_1x_1^2 + a_2x_2^2 - tu^2 = 0$$

$$a_3x_3^2 + a_4x_4^2 + tv^2 = 0.$$

7. Construct the solution of $a_1x_1^2 + a_2x_2^2 + a_3x_3^2 + a_4x_4^2$ from the solutions of two ternary forms.

4.3 Isotropic vector of quadratic form with 5 variables

FUNCTION IsIsotropic5

INPUT: Coefficients a_1, a_2, a_3, a_4, a_5 of a diagonal form $a_1x_1^2 + a_2x_2^2 + a_3x_3^2 + a_4x_4^2 + a_5x_5^2$.

OUTPUT: true if and only if the form is isotropic, and if so, returns an isotropic vector.

Algorithm looks for t such that

$$f_1(\mathbf{x}) = a_1 x_1^2 + a_2 x_2^2 - t u^2$$
 is isotropic and
 $f_2(\mathbf{x}) = a_3 x_3^2 + a_4 x_4^2 + a_5 x_5^2 + t v^2$ is isotropic.

- 1. Solve trivial case, when one of the coefficients is zero.
- 2. Test solvability: coefficients should be not all of the same sign (\sim solvability over R).
- 3. Systematically generate non-trivial pairs (b_1, b_2) of integers. For each pair
 - (a) Compute $t = a_1b_1^2 + a_2b_2^2$.
 - (b) If t = 0 then $(b_1, b_2, 0, 0, 0)$ is isotropic vector.
 - (c) Case $t \neq 0$: solve quaternary quadratic form $a_3x_3^2 + a_4x_4^2 + a_5x_5^2 + tv^2$. If it is isotropic, then construct isotropic vector of the original form from b_1, b_2 and isotropic vector of the quaternary form.

4.4 Isotropic vector of quadratic form with 6 or more variables

FUNCTION IsIsotropic

INPUT: Vector (a_1, a_2, \ldots, a_n) of the coefficients of a diagonal form $a_1x_1^2 + a_2x_2^2 + \cdots + a_nx_n^2$.

OUTPUT: **true** if and only if the form is isotropic, and if so, returns an isotropic vector.

Algorithm looks for t such that

$$f_1(\mathbf{x}) = a_1 x_1^2 + a_2 x_2^2 - t u^2$$
 is isotropic and
 $f_2(\mathbf{x}) = a_3 x_3^2 + \dots + a_n x_n^2 + t v^2$ is isotropic.

- 1. Solve trivial case, when one of the coefficients is zero.
- 2. Test solvability: coefficients should be not all of the same sign (\sim solvability over R). Notice also, whether the coefficients a_3, \ldots, a_n have the same sign it will be helpful when generating t.
- 3. Systematically generate non-trivial pairs (b_1, b_2) of integers. For each pair
 - (a) Compute $t = a_1 b_1^2 + a_2 b_2^2$.
 - (b) If signs of t, a_3, \ldots, a_n are not all the same, then construct an isotropic vector.

(If t = 0 then $(b_1, b_2, 0, 0, ..., 0)$ is isotropic vector. Otherwise solve quadratic form $a_3x_3^2 + \cdots + a_nx_n^2 + tv^2$ and construct isotropic vector afterwards.)

5 Implementation

We implemented the algorithms in Magma. To be able to use it, you need following files:

```
Diophant.spec
QuadForm2.m
QuadForm3.m
nrs.m
QuadForm4.m
QuadForm5.m
```

In Magma, load the package by:

```
> AttachSpec("Diophant.spec");
```

To compute an isotropic vector of the form with integer coefficients a_1, a_2, a_2 (i.e. of the form $a_1x_1^2 + a_2x_2^2 + a_3x_3^2$) type

```
> IsIsotropic3(a1,a2,a3);
```

Functions for the forms with different number of variables work similarly:

```
> IsIsotropic2(a1,a2);
> IsIsotropic4(a1,a2,a3,a4);
> IsIsotropic5(a1,a2,a3,a4,a5);
```

All these functions accept as their argument also a symmetric matrix of particular dimension, for example:

> IsIsotropic3(M);

where M is a symmetric 3 by 3 matrix over integers.

For computing an isotropic vector of the quadratic form with more than 5 variables, there is a function

```
> IsIsotropic(M);
> IsIsotropic(a);
```

accepting a symmetric matrix M or a list **a** of coefficients $a_1, a_2, \ldots a_n$ (in this case the form is of the shape $a_1x_1^2 + a_2x_2^2 + \cdots + a_nx_n^2$).

There is also a test file for the package which can be loaded by

```
> Attach("Test.m");
```

It contains functions

```
> TestForm2(max);
> TestForm3(max);
> TestForm4(max);
> TestForm5(max);
```

taking as an argument a positive integer max, producing random coefficients $a_1, \ldots a_n \in \{-\max, \ldots \max\}$ and solving the form $a_1x_1^2 + \cdots + a_nx_n^2$ for n = 2, 3, 4, 5 (according to the function). These functions after generating and solving each form stop and wait for an input. The input can be **<Enter>** (after that a new form is processed) or q **<Enter>** (after that the functions exit).

For testing the algorithm for the forms of more than 5 variables, there is a similar function

```
> TestForm(n, max);
```

where ${\tt n}$ is the number of variables and ${\tt max}$ determines the range for the coefficients.

For testing the algorithm for the forms given by symmetric matrix there are functions

```
> TestFormM2(max);
```

```
> TestFormM5(max);
```

References

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